

# Multicast

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## Agenda

- Whymulticast?
- Theconceptofmulticast
- Multicastgroups
- Multicastaddressing
- Multicastroutingprotocols
- MBONE
- Multicastapplications
- Conclusions

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## Why multicast?

- Unicast vs. multicast
- typical multicast applications are
  - ◆ audio and
  - ◆ video transmission
  - ◆ collaborating
- multicasting reduces network bandwidth usage

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## Why multicast?

- example of bandwidth usage:
  - ◆ MPEG1 video stream 1.5Mbps
  - ◆ 10 receivers
- bandwidth needed at server site with
  - ◆ unicast 15Mbps
  - ◆ multicast 1.5Mbps

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## The concept of multicast

- IPv4supportsthreetypesof addresses
  - ◆ unicast(point -to-point communication)
  - ◆ broadcast(sendapackettoentire subnet)
  - ◆ multicast(deliveryofpacketstoasetofhoststhathavejoinedtomulticastgroup)

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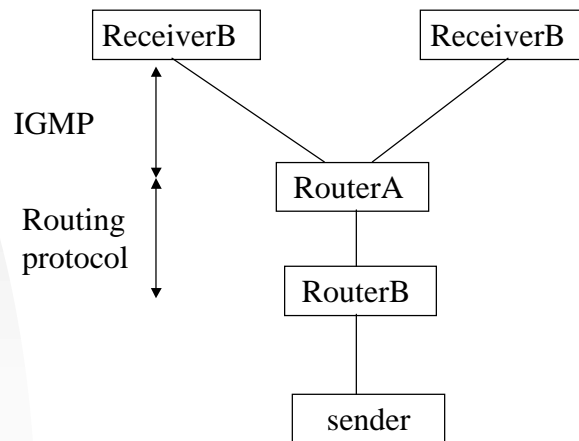
## The concept of multicast

- multicastisnotconnectionoriented
- multicastdatagramisdeliveredtodestinationgroupmembersasaunicastIPdatagram(UDP)
- unicastIPusesclassA,BandC addresses
- multicastIPusesclassDaddress format(224.0.0.0. - 239.255.255.255)

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## The concept of multicast



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## Multicast groups

- hosts are free to join and leave multicast groups
- there are no restrictions on physical locations
- a host may be a member of one or more groups
- a sender need not be a member of the group

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## Multicast addressing

- the format of a 32-bit class D address

1	1	1	0	Multicast Group ID (28 bits)
---	---	---	---	------------------------------

- IANA maintains a list of registered IP multicast groups
- the base address 224.0.0.0 is reserved

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## Multicast addressing

- 224.0.0.1-224.0.0.255 is reserved for routing, discovery and maintenance protocols
  - ◆ 224.0.0.1 all systems on this subnet
  - ◆ 224.0.0.2 all routers on this subnet
  - ◆ 224.0.0.4 all DVMRP routers
  - ◆ 224.0.0.5 all OSPF routers

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## Multicast addressing

- routers should not forward a multicast datagram with a destination address in this range
- groups from 224.0.1.0 to 239.255.255.255 are assigned to various multicast applications or remain unassigned
- 239.0.0.0-239.255.255.255 are reserved for site-local applications, not for Internet-wide applications

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## Multicast addressing

- when sender and receivers are in the same subnet, transmission and reception of multicast frames are simple processes
- sender transmits the IP packet to the multicast group, the NIC maps class D address to network-level multicast address

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## Multicast addressing

- when sender and receivers are located to different subnets, things are more complicated
- routers are required to implement a multicast routing protocol for
  - ◆ discover delivery tree
  - ◆ support multicast data packet forwarding
- each router needs also to implement a group membership protocol to learn group members on its directly attached subnetworks

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## IGMP version 1

- Internet Group Management Protocol (RFC -1112)
- IGMP is used to register a host to multicast group
- routers periodically transmit Host Membership Query messages to determine which groups have members on directly attached networks
- query messages are sent to all - host group (224.0.0.1) with TTL=1

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## IGMP version 1

- when host receives a Query message it responds with Host Membership report to inform which groups it belongs
- routers need not to maintain detailed list which host belong to each multicast group
- routers only need to know that there is at least one group member

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## IGMP version 2

- IGMP2 (RFC -2236) is enhanced version of IGMP1
- backward compatible with IGMP version 1
- Group-Specific Query: Query is sent to specific multicast group rather than all groups
- Leave Group message

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## IGMP version 3

- draft specification
- GroupSourceMessage; host can elect to receive traffic from specific sources of a multicast group
- that will help to reduce bandwidth usage because multicast routing protocols use that information when constructing the multicast delivery trees

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## Multicast forwarding algorithms

- IGMP is used for multicast packet delivery from local router to directly connected subnetworks
- for Internet -wide multicast service there is a need to use multicast routing protocols
- multicast routing protocols are responsible for constructing delivery trees and multicast packet forwarding

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## Multicast forwarding algorithms

- there are different algorithms used by multicast protocols
  - ◆ Flooding
  - ◆ Spanning trees
  - ◆ Reverse Path Broadcasting (RPB)
  - ◆ Truncated Reverse Path Broadcasting (TRPB)
  - ◆ Reverse Path Multicasting (RPM)
  - ◆ Core-Based Trees

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## Flooding

- simplest technique for delivering multicast datagrams
- if the packet is seen first time, router forwards packet to all interfaces
- if router has seen the packet before, it discards packet
- flooding is very simple to implement
- it generates duplicate packets and uses resources inefficiently

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## Spanning trees

- more efficient method than flooding
- spanning tree is a structure of active paths
- packets are forwarded to active paths except to the path which originated the packet
- powerful method and quite easy to implement

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## Reverse Path Broadcasting (RPB)

- source-rooted spanning tree
- router forwards packet to all interfaces (child) except incoming interface (parent)
- distance-vector routing protocol's routing table may be used to determine shortest path back to the source
- need not know whole tree structure

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## Truncated Reverse Path Broadcasting (TRPB)

- RPBforwardspacketstoallleafs
- TRPBeliminatesunnecessary trafficonleafsubnetworksby determininggroupmembershipvia IGMP
- butitdoesnotconsidergroup membershipwhenbuildingthe distributiontree

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## Reverse Path Multicasting (RPM)

- enhancementofRPBAndTRPB
- distributiontreeconsistof
  - ◆ subnetworkswithgroupmembers
  - ◆ shortestpathtosubnetworkswith groupmembers
- prunemechanism;packetsare forwardedonlytobranches that leadtomembersofgroup

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## Reverse Path Multicasting (RPM)

- notscalable;allroutersneedto maintainstateinformationofall groupsandsources
- multicastpacketsmustbe periodicallyforwardedtoevery routerintheinternetworkto maintainpruneinformation

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## Core-Based Trees (CBT)

- RFC-2189
- source-rooted,shortest -path deliverytreeforeachsource, grouppair
- similartothespanningtree algorithmexceptithasdifferent core-basedtreeforeachgroup
- itonlyrequiresaroutertomaintain stateinformationforeachgroup, notforgroup,sourcepair

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## Core-Based Trees (CBT)

- conserve bandwidth since it does not forward multicast packets periodically to every router
- the bottleneck is near a core router because traffic from all sources traverses via same links

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## MOSP

- RFC-1584
- multicast extension to OSPF version 2
- MOSP is defined in RFC -1584
- uses network topology constructed by unicast OSPF
- does not support tunnels

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## MBGP

- MultiprotocolExtensionsforBGP -4 (RFC-2283)
- supportsdifferentunicastand multicastroutingpolicies concurrently
- scalable

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## DVMRP

- DistanceVectorMulticastRouting Protocol(RFC -1075)
- usesReversePathMulticasting protocol
- firstimplementationwasmrouted
- twotypesofinterfaces
  - ◆ physicalinterfaceto directly connectedsubnetworks
  - ◆ tunnelinterfaceto another multicastisland

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## DVMRP

- all interfaces are configured with metric and TTL threshold

Initial TTL Scope	
0	Restricted to the same host
1	Restricted to the same subnet
32	Restricted to the same site
64	Restricted to the same region
128	Restricted to the same continent
255	Unrestricted in scope

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## DVMRP

- multicast packet is forwarded if the TTL value in the IP header is greater than the TTL threshold assigned to the interface
- the initial packet is forwarded to all interfaces to determine prune information
- if there are no members on certain leaf, that branch is removed from delivery tree

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## DVMRP

- if pruned branch discovers new group member it sends graft message to upstream to cancel prune state
- upstream router restores that branch to delivery tree
- DVMRP maintains routing and forwarding tables with integrated routing protocol

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## PIM

- Protocol Independent Multicasting
- PIM is independent on the mechanism provided by any particular unicast routing protocol
- it requires some unicast routing protocol to determine network topology and topology changes
- two modes: dense and sparse

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## PIM-DM

- PIM-DM uses Reverse Path Multicasting
- quite similar to DVMRP
- uses Reverse Path Multicasting (RPM)
- PIM-DM is developed for campus LANs where group membership is usually dense and there is also more bandwidth available than WANs

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## PIM-DM

- differences between PIM and DVMRP
  - ◆ uses existing unicast routing protocol
  - ◆ simply forwards all multicast packets to downstream until prune message is arrived

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## PIM-SM

- RFC-2362
- PIM-SM provides an efficient multicast routing mechanism for communication between sparsely distributed groups
- solves dense mode protocols' scaling problems by limiting multicast packets to routers which are interested in receiving traffic

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## PIM-SM

- PIM-SM differs from existing dense-mode algorithms by following ways
  - ◆ routers are required to join to sparse-mode distribution tree by transmitting join messages
  - ◆ if router is not part of that distribution tree, it will not receive multicast traffic

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## PIM-SM

- PIM-SM uses concept of Rendezvous Points (RP)
- at the RP receivers meet sources
- the initiator of group selects the primary RP and alternative RPs
- only one RP is active
- the host which wants to join the group sends join message to local router which sends explicit join message to group's primary RP

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## MBONE

- Multicast backbone on the Internet
- first transmission was IETF meeting "audiocast" 1992
- MBONE is a virtual network on the top of Internet
- used for audio and video transmissions and for collaborating
- small part of Internet's backbone routers support multicasting

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# MBONE

- typicalsessionsare
  - ◆ IETFmeetings
  - ◆ radiostations
  - ◆ NASAshuttlemissions
  - ◆ eventslikerockconcerts
- typicalbandwidthusage
  - ◆ foraudio13kbps - 64kbps
  - ◆ forvideoupto128kbps

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# MBONE

- mostlyusedbyuniversitiesand researchinstitutes
- mostofthecommercialoperators donotsupportmulticasting
- problemswithfirewalls,one solutionisatunnelinterface
- usuallylowbandwidthtransmission

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## Multicast applications

- at the moment different UNIX systems and also Windows OS support multicasting
- main router vendor support multicasting
- there is free client software available
- the “killer” application is still missing

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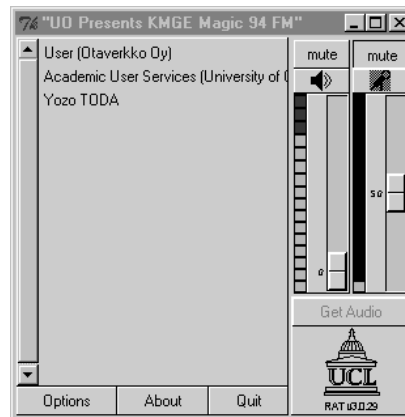
## Multicast applications



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## Multicast applications



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## Multicast applications

- MP3multicasting
  - ◆ offersreasonablequalityforaudio streaming
  - ◆ RTPusedforpayloadformat
  - ◆ mightbe“killerapplication”

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## Conclusions

- multicasting reduces bandwidth usage
- not easy to implement
- needs protocol development work for Internet -wide usage

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## More information

- [www.ietf.org](http://www.ietf.org)
- [www.mbone.com](http://www.mbone.com)

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